Thus
$$|(Surves - Sested bit '10)$$
 there is a deconsistive block-base PIT algorithm for
 $\Xi^{(b)} TL clocks with time complexity poly (ndb). The Sk
day - d.
(Precurves: Kagal-Sareno 1061; where two PIT algorithm) (Thus follows from the follows there in (window reduction).
Thus follows from the follows theorem (window reduction).
Thus 2. (SS'10) let IF be af ske > dink2. There exist maps $d_{1,...,k}$ to $\Xi^{(c)}_{ij}$ if is
with t < puly (kdn) that one computable in these poly (kdn) s.t. for a
 $\Xi^{(b)} T(^{(b)})$ double C, $C = 0 \iff W$ it is $1 > 0$.
Again we assume $C = \Xi^{(c)}_{ij}$ F: where $F_{i} = C^{(c)}_{ij}$ [Li, han genous linear, $C \in F_{i}^{(c)}$
(To construct a bitting-tet H using Thus 2, Con be genoused when $r, C \in F_{i}^{(c)}$
(to H $\in F_{i}^{(c)}$ ca hitting-tet H using Thus 2, Con be genoused when $r \in C_{ij}^{(c)}$. K_{ij}), $K_{ij}^{(c)}$.
Let $H \in F_{i}^{(c)}$ ca hitting-tet H using Thus 2, Con be genoused when $r \in F_{i}^{(c)}$
(to construct a bitting-tet H using Thus 2, Con be genoused when $r \in F_{i}^{(c)}$.
(to H $\in F_{i}^{(c)}$ ca H $= \bigcup P_{i}^{(c)}(H)$ been the Obstandard told
where $f^{(a, \dots, a_{k})} = (\Xi^{(c)}_{i} = i_{i})_{i=1,\dots, n}^{(c)} + j_{i} : r_{i} : J_{i} = J_{i}$.
This a multiplicative term is a polynowid $T = c_{i}^{(c)} I_{i}$. It has some $T_{i} : G$
The radical span of a collection of multiplicative terms $S = \{T_{i}, \dots, T_{r}\}$ is
 $rad sp((S) = rand sp((S_{i}), \dots, T_{r}) = \int \Sigma_{i} : L_{i}$. It dives some $T_{i} : G$
To obser words rading (S) is the line space spanned by linear forms apprecised in S.
 $Tsimple: radisp(X_{i}^{(c)}, K_{i}(S)) X_{ir}) = span(X_{i}, X_{i} + X_{i}, X_{i})$.
(3) Let $C = \frac{T}{T} F_{i}$ where $F_{i} = C_{i} T_{i}^{(c)}$.
 $A path of C of (ength $K' \le k$ is a K -tuple (V_{i}, \dots, V_{k}) of multiplicative terms$$

A park of C of leagth k'sk is a k-taple (V, ..., V,) of multiplicitie terms
such that for each
$$i \in \{1, ..., k'\}$$
, there exists $lin_{i,j} \in approximation = Fill
such that for each $i \in \{1, ..., k'\}$, there exists $lin_{i,j} \in approximation = Fill
such that $V_1 = \prod_{i \in j \leq d} lin_{i,j} = 0$; $i := 1$.
 $lin_{i \in j \leq d} lin_{j}$ and $padip(V_1, ..., V_{-1})$, $c \in F$
In other works, V_i clicks all $lin_{i,j}$ that become a milliple of $lin_{i,j}$ and $radip(V_1, ..., V_{i-1})$.
Example: $(1 : X_i^2 \times X_i X_i + 1 \times X_i + i \times I_i) (X_i + I \times I_i) (X_i + X_i + I \times I_i) (X_i +$$$

$$\begin{array}{l} \forall k \in \mathbb{F} \quad S_{0} \subset = \operatorname{First} + \cdots + \operatorname{Firs} \quad (\operatorname{und} (u_{1}, \dots, v_{1})) \\ & (\neq u, \operatorname{First} + \operatorname{First} + \operatorname{First} \notin (\operatorname{und} (u_{1}, \dots, v_{1})) \\ & (=) \quad F_{1+1} + \cdots + \operatorname{First} \notin (v_{1}, \cdots, v_{1}, \dots, v_{1}) \\ & (=) \quad v_{2} \quad \text{the fact First} + \operatorname{First} \notin (v_{1}, \cdots, v_{1}, \dots, v_{1}) \\ & (=) \quad v_{2} \quad \text{the fact First} + \operatorname{First} \oplus (v_{2}, \dots, v_{2}) \\ & (=) \quad v_{2} \quad \text{operature the olderals.} \\ \\ & (=) \quad F_{1} = \frac{1}{11} \quad \operatorname{First}_{0} \quad \text{where each First}_{1} \quad (j \quad \text{chechs the linest forms thest are each other word matip((v_{1}, \dots, v_{2})) \\ & (=) \quad F_{1} = 1 \quad \text{First}_{1} \quad \text{where each First}_{1} \quad (j \quad \text{chechs other each each first}_{1} \quad (=) \quad F_{1} = 1 \quad F_{1} \in S \quad \text{che error forms}_{1} \quad t_{1} \quad \text{operature each first}_{1} \quad (=) \quad (detacls our) \quad (detacls our) \quad (=) \quad F_{1} = 1 \quad \cdots \quad F_{1} = S \quad (=) \quad (=)$$